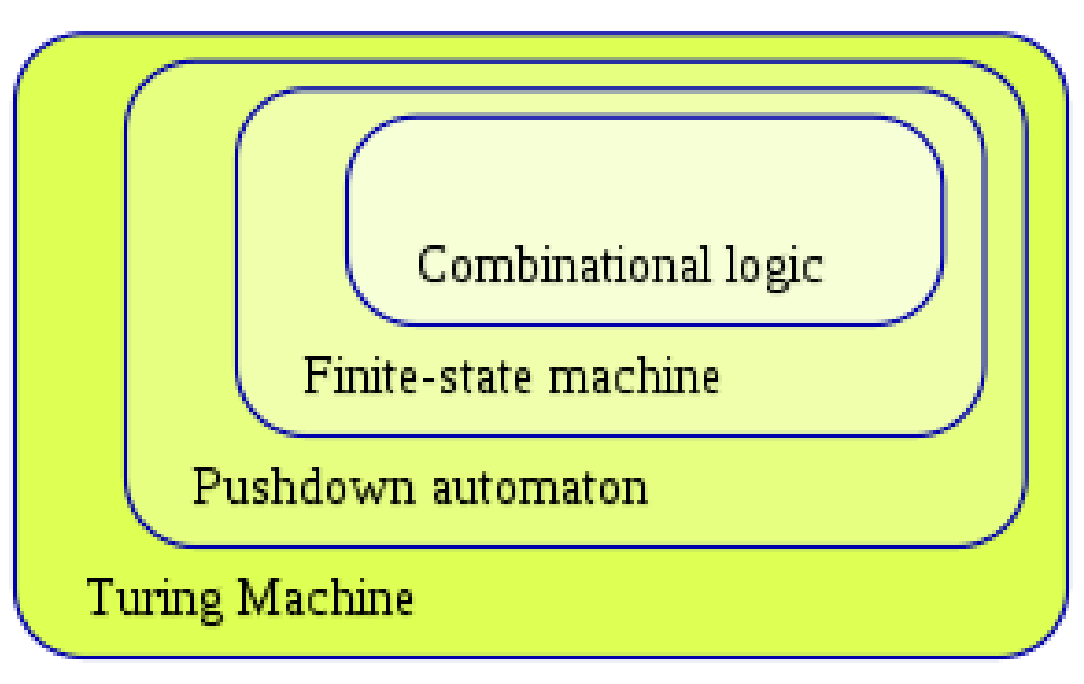


# Computational models



Classes of automata

[https://en.wikipedia.org/wiki/Automata\\_theory](https://en.wikipedia.org/wiki/Automata_theory)

Automaton
Deterministic Finite Automaton (DFA) -- Lowest Power
(same power)    (same power)
Nondeterministic Finite Automaton (NFA)
(above is weaker) ∩ (below is stronger)
Deterministic Push Down Automaton (DPDA-I)
with 1 push-down store
∩
Nondeterministic Push Down Automaton (NPDA-I)
with 1 push-down store
∩
Linear Bounded Automaton (LBA)
∩
Deterministic Push Down Automaton (DPDA-II)
with 2 push-down stores
Nondeterministic Push Down Automaton (NPDA-II)
with 2 push-down stores
Deterministic Turing Machine (DTM)
Nondeterministic Turing Machine (NTM)
Probabilistic Turing Machine (PTM)
Multitape Turing Machine (MTM)
Multidimensional Turing Machine

# (Formal) grammars

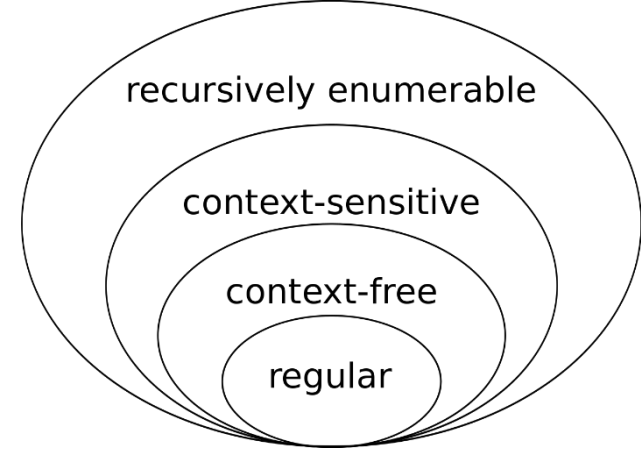
N. Chomsky: separation of syntax and semantics

A grammar is formally defined as the tuple  $(N, \Sigma, P, S)$  transformational grammar/rewrite system

- A finite set  $N$  of *nonterminal symbols*, that is disjoint with the strings formed from  $G$ .
- A finite set  $\Sigma$  of *terminal symbols* that is disjoint from  $N$ .
- A finite set  $P$  of *production rules*, each rule of the form
$$(\Sigma \cup N)^* N (\Sigma \cup N)^* \rightarrow (\Sigma \cup N)^*$$
- A distinguished symbol  $S \in N$  that is the *start symbol*, also called the *sentence symbol*.

[https://en.wikipedia.org/wiki/Formal\\_grammar](https://en.wikipedia.org/wiki/Formal_grammar)

# The Chomsky hierarchy



Grammar	Languages	Automaton	Production rules (constraints)*	Examples <sup>[3]</sup>
Type-0	Recursively enumerable	Turing machine	$\alpha A \beta \rightarrow \gamma$	$L = \{w \mid w \text{ describes a terminating Turing machine}\}$
Type-1	Context-sensitive	Linear-bounded non-deterministic Turing machine	$\alpha A \beta \rightarrow \alpha \gamma \beta$	$L = \{a^n b^n c^n \mid n > 0\}$
Type-2	Context-free	Non-deterministic pushdown automaton	$A \rightarrow \alpha$	$L = \{a^n b^n \mid n > 0\}$
Type-3	Regular	Finite state automaton	$A \rightarrow a$ and $A \rightarrow aB$	$L = \{a^n \mid n \geq 0\}$

\* Meaning of symbols:

- $a$  = terminal
- $A, B$  = non-terminal
- $\alpha, \beta, \gamma$  = string of terminals and/or non-terminals
  - $\alpha, \beta$  = maybe empty
  - $\gamma$  = never empty

[https://en.wikipedia.org/wiki/Chomsky\\_hierarchy](https://en.wikipedia.org/wiki/Chomsky_hierarchy)