# Stochastic inference in Bayesian networks, Markov chain Monte Carlo METHODS 

## AI: Stochastic inference in BNs

## Outline

$\diamond$ Types of inference in (causal) BNs
$\diamond$ Hardness of exact inference in general BNs
$\diamond$ Approximate inference by stochastic simulation
$\diamond$ Approximate inference by Markov chain Monte Carlo

## Inference tasks

Simple queries: compute posterior marginal $\mathbf{P}\left(X_{i} \mid \mathbf{E}=\mathbf{e}\right)$
e.g., $P($ NoGas $\mid$ Gauge $=$ empty, Lights $=o n$, Starts $=$ false $)$

Conjunctive queries: $\mathbf{P}\left(X_{i}, X_{j} \mid \mathbf{E}=\mathbf{e}\right)=\mathbf{P}\left(X_{i} \mid \mathbf{E}=\mathbf{e}\right) \mathbf{P}\left(X_{j} \mid X_{i}, \mathbf{E}=\mathbf{e}\right)$
Optimal decisions: decision networks include utility information; probabilistic inference required for $P$ (outcome|action, evidence)

Value of information: which evidence to seek next?
Sensitivity analysis: which probability values are most critical?
Explanation: why do I need a new starter motor?
Causal inference: what is the effect of an intervention?
Counterfactual inference: what would have been the effect of a hypothetical/imagery past intervention\&observation?

## Inference by enumeration: principle

Let X be all the variables. Typically, we want the posterior joint distribution of the query variables $Y$ given specific values e for the evidence variables $E$.

Let the hidden variables be $\mathrm{H}=\mathrm{X}-\mathrm{Y}-\mathrm{E}$.
Then the required summation of joint entries is done by summing out the hidden variables:

$$
\mathbf{P}(\mathbf{Y} \mid \mathbf{E}=\mathbf{e})=\alpha \mathbf{P}(\mathbf{Y}, \mathbf{E}=\mathbf{e})=\alpha \sum_{\mathbf{h}} \mathbf{P}(\mathbf{Y}, \mathbf{E}=\mathbf{e}, \mathbf{H}=\mathbf{h})
$$

The terms in the summation are joint entries!
Obvious problems:

1) Worst-case time complexity $O\left(d^{n}\right)$ where $d$ is the largest arity
2) Space complexity $O\left(d^{n}\right)$ to store the joint distribution
3) How to find the numbers for $O\left(d^{n}\right)$ entries???

## Complexity of exact inference

Singly connected networks (or polytrees):

- any two nodes are connected by at most one (undirected) path
- time and space cost of exact inference $O\left(d^{k} n\right)$

Multiply connected networks:

- can reduce 3SAT to exact inference: $0<\mathrm{p}($ AND $) ? \Rightarrow$ NP-hard
- equivalent to counting 3SAT models $\Rightarrow$ \#P-complete

1. $A \vee B \vee C$
2. $C \vee D v \neg A$
3. $B \vee C \vee \neg D$


## Inference by stochastic simulation

Basic idea:

1) Draw $N$ samples from a sampling distribution $S$
2) Compute an approximate posterior probability
3) Show this converges to the true probability $P$

Outline:

- Sampling from an empty network
- Rejection sampling: reject samples disagreeing with evidence
- Likelihood weighting: use evidence to weight samples
- Markov chain Monte Carlo (MCMC): sample from a stochastic process
whose stationary distribution is the true posterior


## Sampling from an empty network

```
function PRIOR-SAMPLE( }bn\mathrm{ ) returns an event sampled from bn
    inputs: bn, a belief network specifying joint distribution }\mathbf{P}(\mp@subsup{X}{1}{},\ldots,\mp@subsup{X}{n}{}
    x}\leftarrow\mathrm{ an event with }n\mathrm{ elements
    for }i=1\mathrm{ to }n\mathrm{ do
        x}\leftarrow\mp@code{a random sample from }\mathbf{P}(\mp@subsup{X}{i}{}|\operatorname{parents}(\mp@subsup{X}{i}{})
        given the values of Parents( }\mp@subsup{X}{i}{})\mathrm{ in x
    return x
```

Example

Example

Example

Example

Example




## Rejection sampling

$\hat{\mathbf{P}}(X \mid \mathbf{e})$ estimated from samples agreeing with e

```
function REJECTION-SAMPLING \((X, \mathbf{e}, b n, N)\) returns an estimate of \(P(X \| \mathbf{e})\)
    local variables: \(\mathbf{N}\), a vector of counts over \(X\), initially zero
    for \(j=1\) to \(N\) do
        \(\mathbf{x} \leftarrow \operatorname{PRIOR}-\operatorname{SAMPLE}(b n)\)
        if x is consistent with e then
            \(\mathbf{N}[x] \leftarrow \mathbf{N}[x]+1\) where \(x\) is the value of \(X\) in \(\mathbf{x}\)
    return Normalize( \(\mathrm{N}[X]\) )
```

E.g., estimate $\mathbf{P}($ Rain $\mid$ Sprinkler $=$ true $)$ using 100 samples
27 samples have Sprinkler $=$ true
Of these, 8 have Rain=true and 19 have Rain=false.
$\hat{\mathbf{P}}($ Rain $\mid$ Sprinkler $=$ true $)=\operatorname{NORMALIZE}(\langle 8,19\rangle)=\langle 0.296,0.704\rangle$

## Analysis of rejection sampling

$$
\begin{array}{ll}
\hat{\mathbf{P}}(X \mid \mathbf{e})=\alpha \mathbf{N}_{P S}(X, \mathbf{e}) \quad \text { (algorithm defn.) } \\
& \left.=\mathbf{N}_{P S}(X, \mathbf{e}) / N_{P S}(\mathbf{e}) \quad \text { (normalized by } N_{P S}(\mathbf{e})\right) \\
& \approx \mathbf{P}(X, \mathbf{e}) / P(\mathbf{e}) \quad \text { (property of PrIORSAMPLE) } \\
& =\mathbf{P}(X \mid \mathbf{e}) \quad \text { (defn. of conditional probability) }
\end{array}
$$

Hence rejection sampling returns consistent posterior estimates
Problem: hopelessly expensive if $P(\mathbf{e})$ is small
$P($ e $)$ drops off exponentially with number of evidence variables!

## The Markov chain

With Sprinkler $=$ true, WetGrass $=$ true, there are four states:


Wander about for a while, average what you see

## Markov blanket

Each node is conditionally independent of all others given its Markov blanket: parents + children + children's parents


## Approximate inference using MCMC

"State" of network = current assignment to all variables. Generate next state by sampling one variable given Markov blanket. Sample each variable in turn, keeping evidence fixed

```
function MCMC- \(\operatorname{Ask}(X, \mathbf{e}, b n, N)\) returns an estimate of \(P(X \mid \mathbf{e})\)
    local variables: \(\mathbf{N}[X]\), a vector of counts over \(X\), initially zero
                                    \(\mathbf{Z}\), the nonevidence variables in \(b n\)
                                    x , the current state of the network, initially copied from
```

    initialize x with random values for the variables in Y
    for \(j=1\) to \(N\) do
        for each \(Z_{i}\) in Z do
            sample the value of \(Z_{i}\) in \(\mathbf{x}\) from \(\mathbf{P}\left(Z_{i} \mid m b\left(Z_{i}\right)\right)\)
                given the values of \(M B\left(Z_{i}\right)\) in \(\mathbf{x}\)
            \(\mathbf{N}[x] \leftarrow \mathbf{N}[x]+1\) where \(x\) is the value of \(X\) in \(\mathbf{x}\)
    return Normalize( \(\mathrm{N}[X]\) )
    Can also choose a variable to sample at random each time

## MCMC example contd.

## Estimate $\mathbf{P}($ Rain $\mid$ Sprinkler $=$ true, WetGrass $=$ true $)$

Sample Cloudy or Rain given its Markov blanket, repeat. Count number of times Rain is true and false in the samples.
E.g., visit 100 states

31 have Rain = true, 69 have Rain $=$ false
$\hat{\mathbf{P}}($ Rain $\mid$ Sprinkler $=$ true, WetGrass $=$ true $)$
$=$ Normalize $(\langle 31,69\rangle)=\langle 0.31,0.69\rangle$
Theorem: chain approaches stationary distribution: long-run fraction of time spent in each state is exactly proportional to its posterior probability

## Markov blanket sampling

Markov blanket of Cloudy is Sprinkler and Rain
Markov blanket of Rain is
Cloudy, Sprinkler, and WetGrass


Probability given the Markov blanket is calculated as follows:

$$
P\left(x_{i}^{\prime} \mid m b\left(X_{i}\right)\right)=P\left(x_{i}^{\prime} \mid \operatorname{parents}\left(X_{i}\right)\right) \Pi_{Z_{j} \in \operatorname{Children}\left(X_{i}\right)} P\left(z_{j} \mid \text { parents }\left(Z_{j}\right)\right)
$$

Easily implemented in message-passing parallel systems, brains
Main computational problems:

1) Difficult to tell if convergence has been achieved
2) Can be wasteful if Markov blanket is large: $P\left(X_{i} \mid m b\left(X_{i}\right)\right)$ won't change much (law of large numbers)

$\diamond$ Hardware!
$\diamond$ Bayesian model averaging

## MCMC analysis: Outline

Transition probability $q\left(\mathrm{x} \rightarrow \mathrm{x}^{\prime}\right)$
Occupancy probability $\pi_{t}(\mathbf{x})$ at time $t$
Equilibrium condition on $\pi_{t}$ defines stationary distribution $\pi(\mathbf{x})$
Note: stationary distribution depends on choice of $q\left(\mathrm{x} \rightarrow \mathrm{x}^{\prime}\right)$
Pairwise detailed balance on states guarantees equilibrium
Gibbs sampling transition probability:
sample each variable given current values of all others
$\Rightarrow$ detailed balance with the true posterior
For Bayesian networks, Gibbs sampling reduces to
sampling conditioned on each variable's Markov blanket

## Stationary distribution

$\pi_{t}(\mathrm{x})=$ probability in state x at time $t$
$\pi_{t+1}\left(\mathrm{x}^{\prime}\right)=$ probability in state $\mathrm{x}^{\prime}$ at time $t+1$
$\pi_{t+1}$ in terms of $\pi_{t}$ and $q\left(\mathbf{x} \rightarrow \mathrm{x}^{\prime}\right)$

$$
\pi_{t+1}\left(\mathbf{x}^{\prime}\right)=\sum_{\mathbf{x}} \pi_{t}(\mathbf{x}) q\left(\mathbf{x} \rightarrow \mathbf{x}^{\prime}\right)
$$

Stationary distribution: $\pi_{t}=\pi_{t+1}=\pi$

$$
\pi\left(\mathbf{x}^{\prime}\right)=\sum_{\mathbf{x}} \pi(\mathbf{x}) q\left(\mathbf{x} \rightarrow \mathbf{x}^{\prime}\right) \quad \text { for all } \mathbf{x}^{\prime}
$$

If $\pi$ exists, it is unique (specific to $q\left(\mathbf{x} \rightarrow \mathrm{x}^{\prime}\right)$ )
In equilibrium, expected "outflow" = expected "inflow"

## Detailed balance

"Outflow" = "inflow" for each pair of states:

$$
\pi(\mathbf{x}) q\left(\mathbf{x} \rightarrow \mathbf{x}^{\prime}\right)=\pi\left(\mathbf{x}^{\prime}\right) q\left(\mathbf{x}^{\prime} \rightarrow \mathbf{x}\right) \quad \text { for all } \mathbf{x}, \mathbf{x}^{\prime}
$$

Detailed balance $\Rightarrow$ stationarity:

$$
\begin{aligned}
\sum_{\mathbf{x}} \pi(\mathbf{x}) q\left(\mathbf{x} \rightarrow \mathbf{x}^{\prime}\right) & =\sum_{\mathbf{x}} \pi\left(\mathbf{x}^{\prime}\right) q\left(\mathbf{x}^{\prime} \rightarrow \mathbf{x}\right) \\
& =\pi\left(\mathbf{x}^{\prime}\right) \sum_{\mathbf{x}} q\left(\mathbf{x}^{\prime} \rightarrow \mathbf{x}\right) \\
& =\pi\left(\mathbf{x}^{\prime}\right)
\end{aligned}
$$

MCMC algorithms typically constructed by designing a transition probability $q$ that is in detailed balance with desired $\pi$

## Gibbs sampling

Sample each variable in turn, given all other variables
Sampling $X_{i}$, let $\overline{\mathbf{X}}_{i}$ be all other nonevidence variables Current values are $x_{i}$ and $\overline{\mathbf{x}}_{i} ; \mathrm{e}$ is fixed Transition probability is given by

$$
q\left(\mathbf{x} \rightarrow \mathbf{x}^{\prime}\right)=q\left(x_{i}, \overline{\mathbf{x}}_{i} \rightarrow x_{i}^{\prime}, \overline{\mathbf{x}}_{i}\right)=P\left(x_{i}^{\prime} \mid \overline{\mathbf{x}}_{i}, \mathbf{e}\right)
$$

This gives detailed balance with true posterior $P(\mathrm{x} \mid \mathrm{e})$ :

$$
\begin{aligned}
\pi(\mathbf{x}) q\left(\mathbf{x} \rightarrow \mathbf{x}^{\prime}\right) & =P(\mathbf{x} \mid \mathbf{e}) P\left(x_{i}^{\prime} \mid \overline{\mathbf{x}}_{i}, \mathbf{e}\right)=P\left(x_{i}, \overline{\mathbf{x}}_{i} \mid \mathbf{e}\right) P\left(x_{i}^{\prime} \mid \overline{\mathbf{x}}_{i}, \mathbf{e}\right) \\
& =P\left(x_{i} \mid \overline{\mathbf{x}}_{i}, \mathbf{e}\right) P\left(\overline{\mathbf{x}}_{i} \mid \mathbf{e}\right) P\left(x_{i}^{\prime} \mid \overline{\mathbf{x}}_{i}, \mathbf{e}\right) \quad \text { (chain rule) } \\
& =P\left(x_{i} \mid \overline{\mathbf{x}}_{i}, \mathbf{e}\right) P\left(x_{i}^{\prime}, \overline{\mathbf{x}}_{i} \mid \mathbf{e}\right) \quad \text { (chain rule backwards) } \\
& =q\left(\mathbf{x}^{\prime} \rightarrow \mathbf{x}\right) \pi\left(\mathbf{x}^{\prime}\right)=\pi\left(\mathbf{x}^{\prime}\right) q\left(\mathbf{x}^{\prime} \rightarrow \mathbf{x}\right)
\end{aligned}
$$

## Performance of approximation algorithms

Absolute approximation: $|P(X \mid \mathrm{e})-\hat{P}(X \mid \mathrm{e})| \leq \epsilon$
Relative approximation: $\frac{|P(X \mid \mathbf{e})-\hat{P}(X \mid \mathbf{e})|}{P(X \mid \mathbf{e})} \leq \epsilon$
Relative $\Rightarrow$ absolute since $0 \leq P \leq 1$ (may be $O\left(2^{-n}\right)$ )
Randomized algorithms may fail with probability at most $\delta$
Polytime approximation: poly $\left(n, \epsilon^{-1}, \log \delta^{-1}\right)$
Theorem (Dagum and Luby, 1993): both absolute and relative approximation for either deterministic or randomized algorithms are NP-hard for any $\epsilon, \delta<0.5$
(Absolute approximation polytime with no evidence-Chernoff bounds)

## Summary

Exact inference:

- polytime on polytrees (NBNs,HMMs), NP-hard on general graphs
- space $=$ time, very sensitive to topology

Approximate inference:

- Convergence can be very slow with probabilities close to 1 or 0
- Can handle arbitrary combinations of discrete and continuous variables

