# STOCHASTIC INFERENCE IN BAYESIAN NETWORKS, MARKOV CHAIN MONTE CARLO METHODS

#### AI: STOCHASTIC INFERENCE IN BNS

# Outline

- $\diamondsuit$  Types of inference in (causal) BNs
- $\diamondsuit$  Hardness of exact inference in general BNs
- $\diamondsuit$  Approximate inference by stochastic simulation
- $\diamondsuit$  Approximate inference by Markov chain Monte Carlo

### Inference tasks

Simple queries: compute posterior marginal  $\mathbf{P}(X_i | \mathbf{E} = \mathbf{e})$ e.g., P(NoGas | Gauge = empty, Lights = on, Starts = false)

Conjunctive queries:  $\mathbf{P}(X_i, X_j | \mathbf{E} = \mathbf{e}) = \mathbf{P}(X_i | \mathbf{E} = \mathbf{e})\mathbf{P}(X_j | X_i, \mathbf{E} = \mathbf{e})$ 

Optimal decisions: decision networks include utility information; probabilistic inference required for P(outcome|action, evidence)

Value of information: which evidence to seek next?

Sensitivity analysis: which probability values are most critical?

Explanation: why do I need a new starter motor?

Causal inference: what is the effect of an intervention?

Counterfactual inference: what would have been the effect of a hypothetical/imagery past intervention&observation?

## Inference by enumeration: principle

Let X be all the variables. Typically, we want the posterior joint distribution of the query variables Y given specific values e for the evidence variables E.

Let the hidden variables be H = X - Y - E.

Then the required summation of joint entries is done by summing out the hidden variables:

 $\mathbf{P}(\mathbf{Y}|\mathbf{E} = \mathbf{e}) = \alpha \mathbf{P}(\mathbf{Y}, \mathbf{E} = \mathbf{e}) = \alpha \Sigma_{\mathbf{h}} \mathbf{P}(\mathbf{Y}, \mathbf{E} = \mathbf{e}, \mathbf{H} = \mathbf{h})$ 

The terms in the summation are joint entries!

Obvious problems:

- 1) Worst-case time complexity  $O(d^n)$  where d is the largest arity
- 2) Space complexity  $O(d^n)$  to store the joint distribution
- 3) How to find the numbers for  $O(d^n)$  entries???

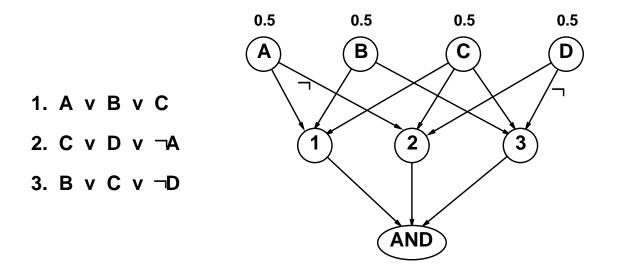
### **Complexity of exact inference**

Singly connected networks (or polytrees):

- any two nodes are connected by at most one (undirected) path
- time and space cost of exact inference  $O(d^k n)$

Multiply connected networks:

- can reduce 3SAT to exact inference:  $0 < p(AND)? \Rightarrow NP$ -hard
- equivalent to **counting** 3SAT models  $\Rightarrow$  #P-complete



# Inference by stochastic simulation

Basic idea:

- 1) Draw N samples from a sampling distribution S
- 2) Compute an approximate posterior probability  $\hat{P}$
- 3) Show this converges to the true probability P

Outline:

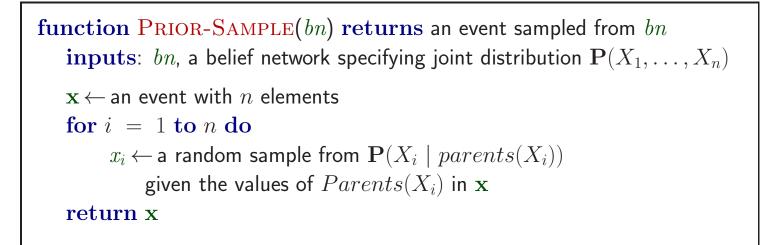
- Sampling from an empty network
- Rejection sampling: reject samples disagreeing with evidence
- Likelihood weighting: use evidence to weight samples
- Markov chain Monte Carlo (MCMC): sample from a stochastic

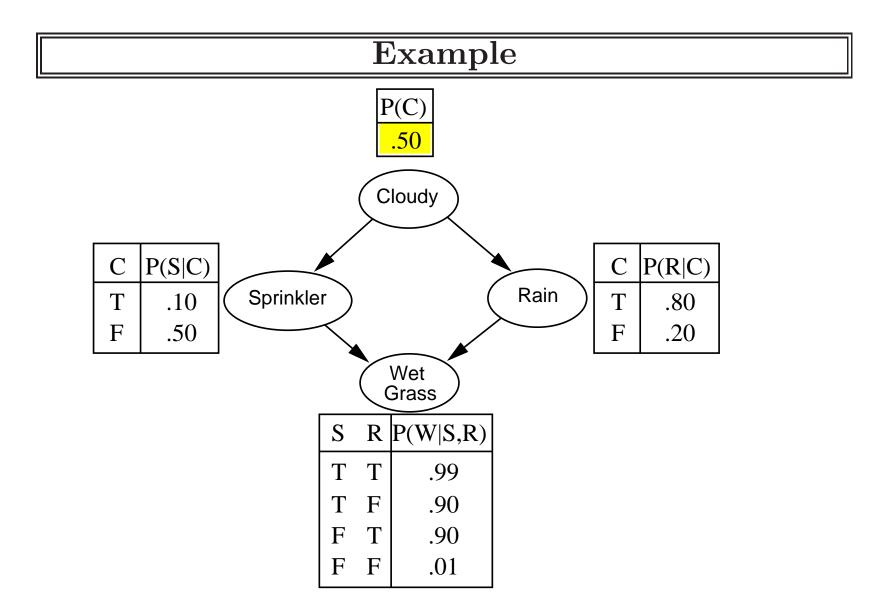
process

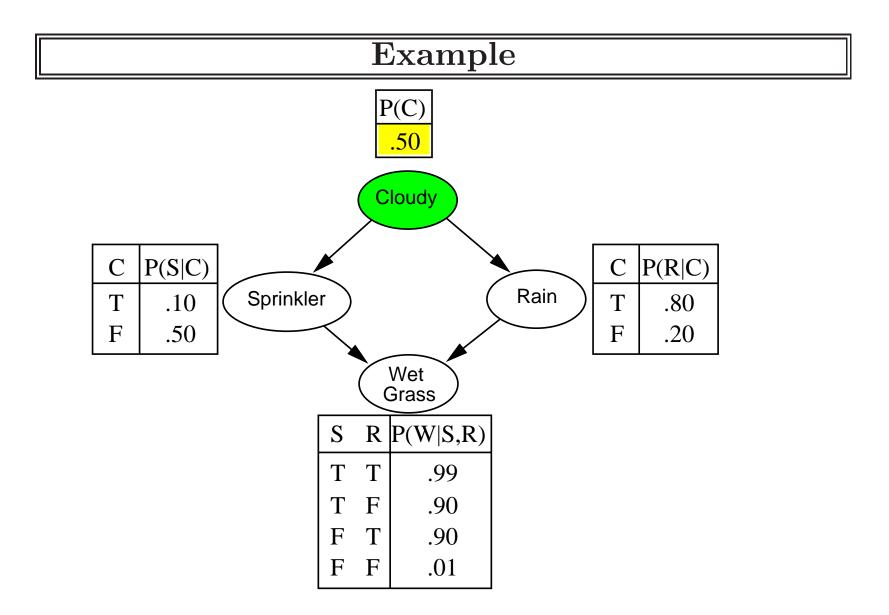
whose stationary distribution is the true posterior

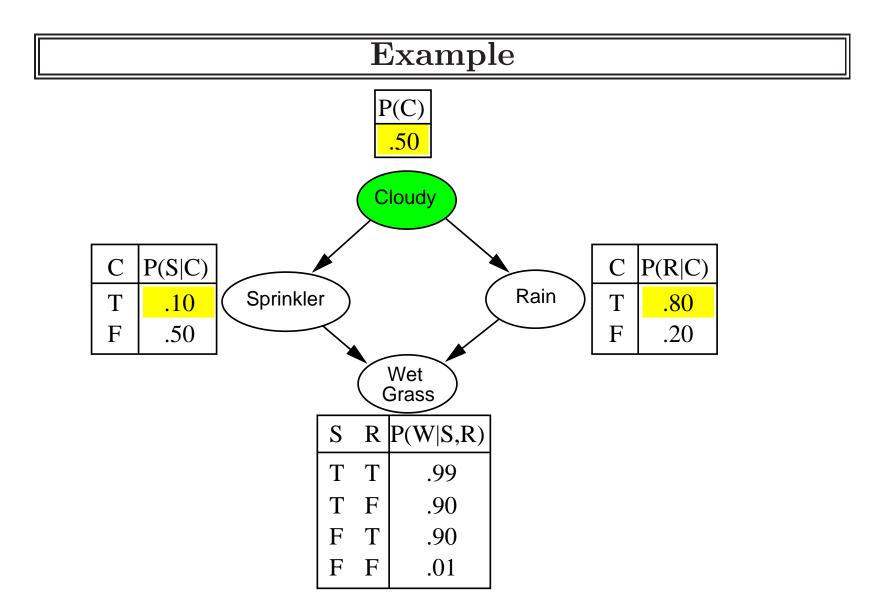


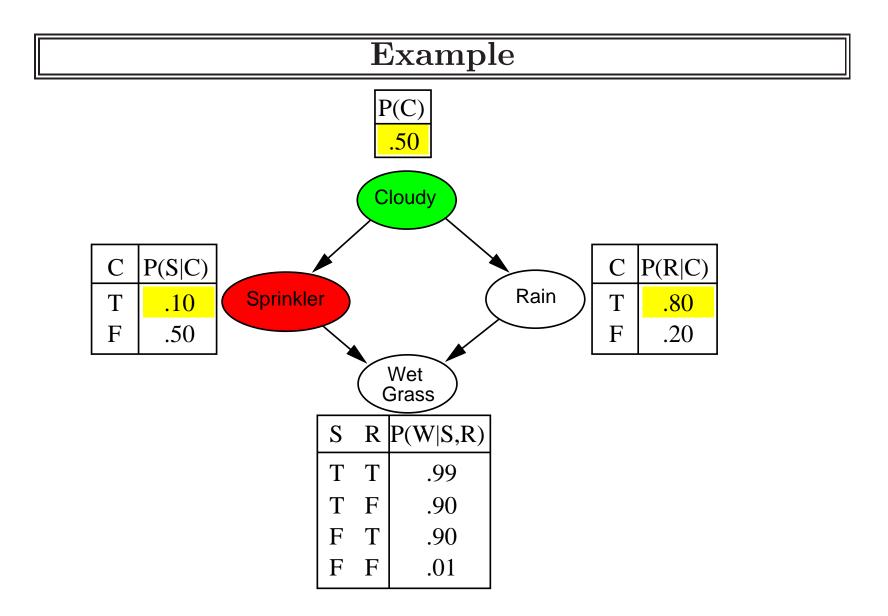
### Sampling from an empty network

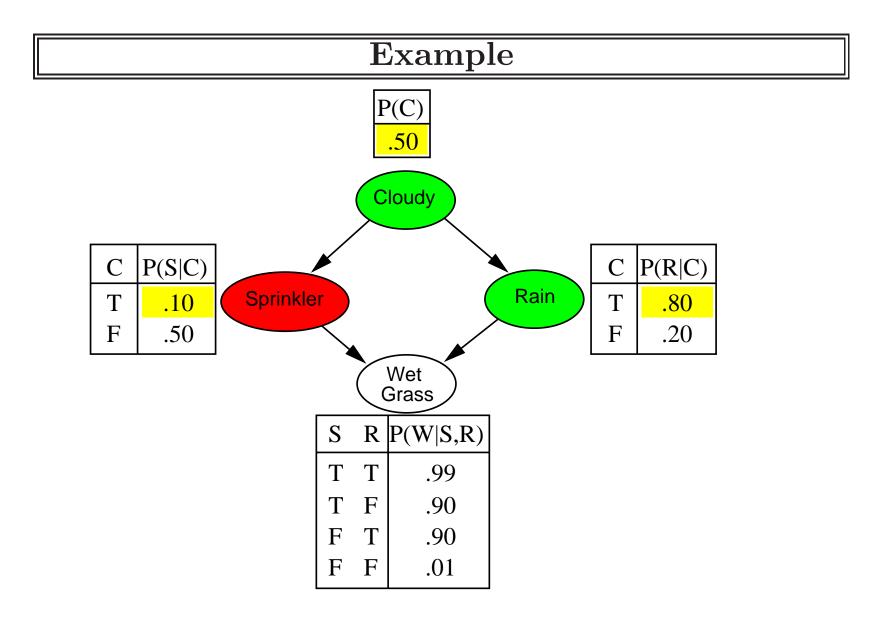


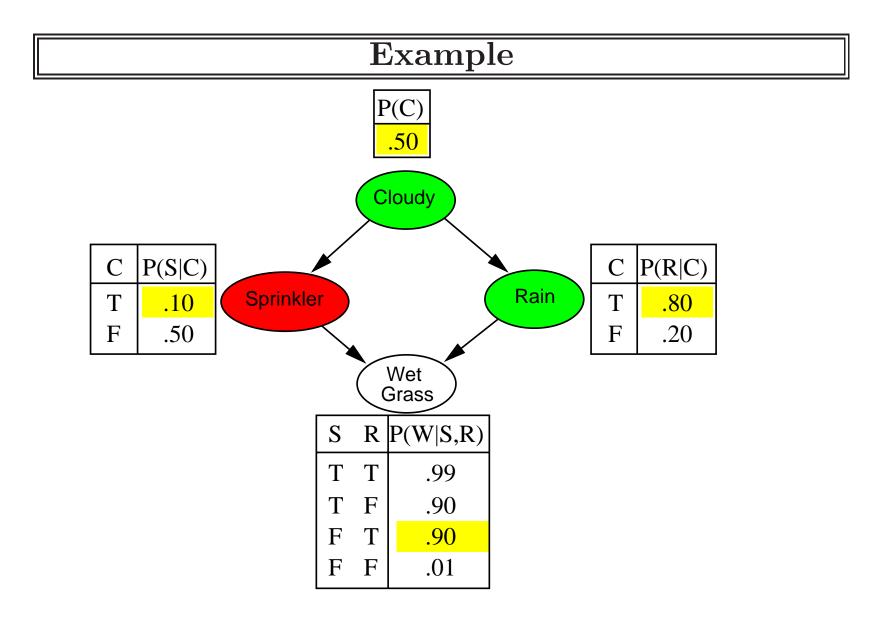


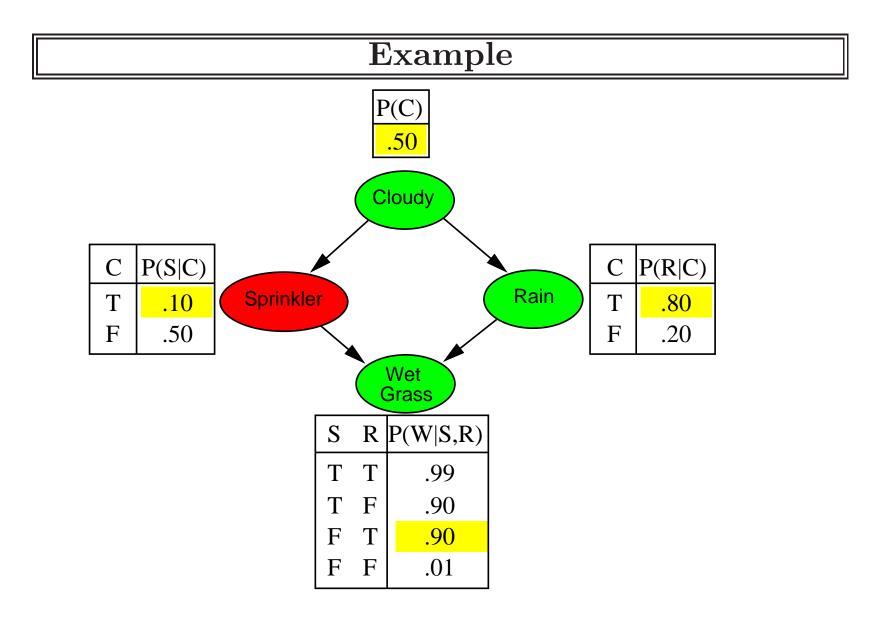












# **Rejection** sampling

 $\hat{\mathbf{P}}(X|\mathbf{e})$  estimated from samples agreeing with  $\mathbf{e}$ 

function REJECTION-SAMPLING(X, e, bn, N) returns an estimate of P(X | e)local variables: N, a vector of counts over X, initially zero for j = 1 to N do  $x \leftarrow PRIOR-SAMPLE(bn)$ if x is consistent with e then  $N[x] \leftarrow N[x]+1$  where x is the value of X in x return NORMALIZE(N[X])

E.g., estimate  $\mathbf{P}(Rain|Sprinkler = true)$  using 100 samples 27 samples have Sprinkler = trueOf these, 8 have Rain = true and 19 have Rain = false.

 $\hat{\mathbf{P}}(Rain|Sprinkler = true) = \text{NORMALIZE}(\langle 8, 19 \rangle) = \langle 0.296, 0.704 \rangle$ 

# Analysis of rejection sampling

$$\begin{split} \hat{\mathbf{P}}(X|\mathbf{e}) &= \alpha \mathbf{N}_{PS}(X,\mathbf{e}) & \text{(algorithm defn.)} \\ &= \mathbf{N}_{PS}(X,\mathbf{e})/N_{PS}(\mathbf{e}) & \text{(normalized by } N_{PS}(\mathbf{e})) \\ &\approx \mathbf{P}(X,\mathbf{e})/P(\mathbf{e}) & \text{(property of PRIORSAMPLE)} \\ &= \mathbf{P}(X|\mathbf{e}) & \text{(defn. of conditional probability)} \end{split}$$

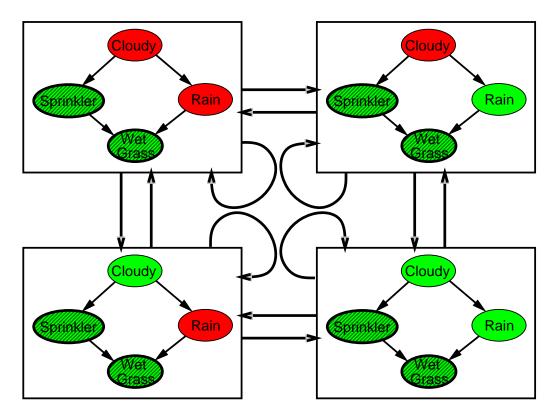
Hence rejection sampling returns consistent posterior estimates

Problem: hopelessly expensive if  $P(\mathbf{e})$  is small

 $P(\mathbf{e})$  drops off exponentially with number of evidence variables!

### The Markov chain

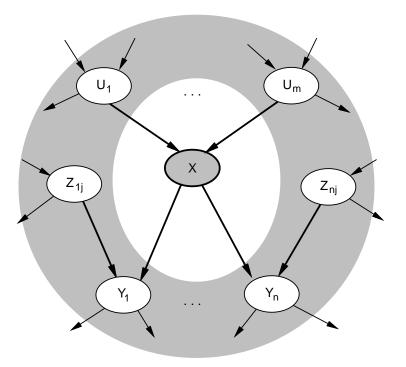
With Sprinkler = true, WetGrass = true, there are four states:



Wander about for a while, average what you see

### Markov blanket

Each node is conditionally independent of all others given its Markov blanket: parents + children + children's parents



# Approximate inference using MCMC

"State" of network = current assignment to all variables. Generate next state by sampling one variable given Markov blanket. Sample each variable in turn, keeping evidence fixed

```
function MCMC-Ask(X, e, bn, N) returns an estimate of P(X|e)
local variables: N[X], a vector of counts over X, initially zero
Z, the nonevidence variables in bn
x, the current state of the network, initially copied from e
initialize x with random values for the variables in Y
for j = 1 to N do
for each Z_i in Z do
sample the value of Z_i in x from P(Z_i|mb(Z_i))
given the values of MB(Z_i) in x
N[x] \leftarrow N[x] + 1 where x is the value of X in x
return NORMALIZE(N[X])
```

Can also choose a variable to sample at random each time

# MCMC example contd.

 $\textbf{Estimate } \mathbf{P}(Rain|Sprinkler = true, WetGrass = true)$ 

Sample *Cloudy* or *Rain* given its Markov blanket, repeat. Count number of times *Rain* is true and false in the samples.

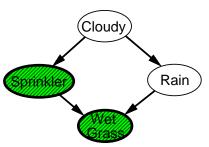
E.g., visit 100 states 31 have Rain = true, 69 have Rain = false

 $\hat{\mathbf{P}}(Rain|Sprinkler = true, WetGrass = true) = \text{NORMALIZE}(\langle 31, 69 \rangle) = \langle 0.31, 0.69 \rangle$ 

Theorem: chain approaches stationary distribution: long-run fraction of time spent in each state is exactly proportional to its posterior probability

# Markov blanket sampling

Markov blanket of *Cloudy* is *Sprinkler* and *Rain* Markov blanket of *Rain* is *Cloudy, Sprinkler*, and *WetGrass* 



Probability given the Markov blanket is calculated as follows:  $P(x'_i|mb(X_i)) = P(x'_i|parents(X_i))\prod_{Z_j \in Children(X_i)} P(z_j|parents(Z_j))$ 

Easily implemented in message-passing parallel systems, brains

Main computational problems:

- 1) Difficult to tell if convergence has been achieved
- 2) Can be wasteful if Markov blanket is large:

 $P(X_i|mb(X_i))$  won't change much (law of large numbers)

# The MCMC age

#### $\diamond$ Hardware!

 $\diamondsuit$  Bayesian model averaging

# MCMC analysis: Outline

Transition probability  $q(\mathbf{x} \rightarrow \mathbf{x'})$ 

Occupancy probability  $\pi_t(\mathbf{x})$  at time t

Equilibrium condition on  $\pi_t$  defines stationary distribution  $\pi(\mathbf{x})$ Note: stationary distribution depends on choice of  $q(\mathbf{x} \to \mathbf{x'})$ 

Pairwise detailed balance on states guarantees equilibrium

Gibbs sampling transition probability:

sample each variable given current values of all others

 $\Rightarrow$  detailed balance with the true posterior

For Bayesian networks, Gibbs sampling reduces to sampling conditioned on each variable's Markov blanket

### **Stationary distribution**

 $\pi_t(\mathbf{x}) = \text{probability in state } \mathbf{x} \text{ at time } t$  $\pi_{t+1}(\mathbf{x}') = \text{probability in state } \mathbf{x}' \text{ at time } t + 1$ 

 $\pi_{t+1}$  in terms of  $\pi_t$  and  $q(\mathbf{x} \to \mathbf{x'})$ 

 $\pi_{t+1}(\mathbf{x}') = \sum_{\mathbf{x}} \pi_t(\mathbf{x}) q(\mathbf{x} \to \mathbf{x}')$ 

Stationary distribution:  $\pi_t = \pi_{t+1} = \pi$ 

$$\pi(\mathbf{x}') = \Sigma_{\mathbf{x}} \pi(\mathbf{x}) q(\mathbf{x} \to \mathbf{x}') \qquad \text{for all } \mathbf{x}'$$

If  $\pi$  exists, it is unique (specific to  $q(\mathbf{x} \rightarrow \mathbf{x}'))$ 

In equilibrium, expected "outflow" = expected "inflow"

### **Detailed balance**

"Outflow" = "inflow" for each pair of states:

 $\pi(\mathbf{x})q(\mathbf{x}\to\mathbf{x}')=\pi(\mathbf{x}')q(\mathbf{x}'\to\mathbf{x})\qquad\text{for all }\mathbf{x},\ \mathbf{x}'$ 

Detailed balance  $\Rightarrow$  stationarity:

$$\Sigma_{\mathbf{x}} \pi(\mathbf{x}) q(\mathbf{x} \to \mathbf{x}') = \Sigma_{\mathbf{x}} \pi(\mathbf{x}') q(\mathbf{x}' \to \mathbf{x})$$
$$= \pi(\mathbf{x}') \Sigma_{\mathbf{x}} q(\mathbf{x}' \to \mathbf{x})$$
$$= \pi(\mathbf{x}')$$

MCMC algorithms typically constructed by designing a transition probability q that is in detailed balance with desired  $\pi$ 

# Gibbs sampling

Sample each variable in turn, given all other variables

Sampling  $X_i$ , let  $\overline{\mathbf{X}}_i$  be all other nonevidence variables Current values are  $x_i$  and  $\overline{\mathbf{x}}_i$ ; **e** is fixed Transition probability is given by

 $q(\mathbf{x} \to \mathbf{x}') = q(x_i, \bar{\mathbf{x}}_i \to x'_i, \bar{\mathbf{x}}_i) = P(x'_i | \bar{\mathbf{x}}_i, \mathbf{e})$ 

This gives detailed balance with true posterior  $P(\mathbf{x}|\mathbf{e})$ :

$$\pi(\mathbf{x})q(\mathbf{x} \to \mathbf{x}') = P(\mathbf{x}|\mathbf{e})P(x'_i|\bar{\mathbf{x}}_i, \mathbf{e}) = P(x_i, \bar{\mathbf{x}}_i|\mathbf{e})P(x'_i|\bar{\mathbf{x}}_i, \mathbf{e})$$
  
=  $P(x_i|\bar{\mathbf{x}}_i, \mathbf{e})P(\bar{\mathbf{x}}_i|\mathbf{e})P(x'_i|\bar{\mathbf{x}}_i, \mathbf{e})$  (chain rule)  
=  $P(x_i|\bar{\mathbf{x}}_i, \mathbf{e})P(x'_i, \bar{\mathbf{x}}_i|\mathbf{e})$  (chain rule backwards)  
=  $q(\mathbf{x}' \to \mathbf{x})\pi(\mathbf{x}') = \pi(\mathbf{x}')q(\mathbf{x}' \to \mathbf{x})$ 

# Performance of approximation algorithms

Absolute approximation:  $|P(X|\mathbf{e}) - \hat{P}(X|\mathbf{e})| \leq \epsilon$ 

Relative approximation:  $\frac{|P(X|\mathbf{e}) - \hat{P}(X|\mathbf{e})|}{P(X|\mathbf{e})} \leq \epsilon$ 

Relative  $\Rightarrow$  absolute since  $0 \le P \le 1 \pmod{(may be O(2^{-n}))}$ 

Randomized algorithms may fail with probability at most  $\delta$ 

Polytime approximation:  $\operatorname{poly}(n, \epsilon^{-1}, \log \delta^{-1})$ 

Theorem (Dagum and Luby, 1993): both absolute and relative approximation for either deterministic or randomized algorithms are NP-hard for any  $\epsilon, \delta < 0.5$ 

(Absolute approximation polytime with no evidence—Chernoff bounds)

#### Summary

Exact inference:

- polytime on polytrees (NBNs,HMMs), NP-hard on general graphs

- space = time, very sensitive to topology

Approximate inference:

– Convergence can be very slow with probabilities close to 1 or 0

- Can handle arbitrary combinations of discrete and continuous variables